

Modeling questions and responses

Lecture 4: the dynamics of responses

Kyle Rawlins

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Johns Hopkins University, Department of Cognitive Science

Outline

Some empirical desiderata

Stalnakerian context

 Responding to assertions

A Stalnakerian account for questioning

Questions and the table

 Polar questions as the tip of the iceberg

Course structure

- *Lecture 1*: Introducing questions and responses.
- *Lecture 2*: Representing question meanings.
- *Lecture 3*: The architecture of a QA system.
- ⇒ *Lecture 4-5*: The dynamics of responses.
- *Lecture 5*: wrap-up.

Some empirical desiderata

Responses

Reminder: the class of responses is large, and answers proper are only a small piece of the picture.

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Goal of Thursday/Friday: give a thorough linguistic account of the pragmatics of responding that derives some of this larger picture.

- **Update semantics with tables.** (Partly based on joint work with Justin Bledin, though he hasn't seen this version.)
- We won't include everything there is, but the account will be flexible, and could be added to.

Responses to assertions

- (1) A: It's raining.
B: I agree.
B': No it's not, that's snow.
B'': Are you sure?
B''': I think there's a water leak on the top floor.

Responses to questions

(2) A: Is it raining?

B: Yes, it is. / No, it isn't.

B': It might be.

B'': I don't know.

B''': I refuse to answer. / fuck you! / (shushing motion)

Responses to questions

- (2) A: Is it raining?
B: Yes, it is. / No, it isn't.
B': It might be.
B'': I don't know.
B''': I refuse to answer. / fuck you! / (shushing motion)
- (3) A: When's the poster session today?
B: It's at 8.
B'/A: Is it in the evening?
B'': There's no poster session today.
B''': It might be at 8.
B''': I don't know.

Stalnakerian context

Stalnakerian context and assertion

- (4) A **context set** is a set of worlds. (Stalnaker 1978)
- (5) **Contexts v. 1**: A context is a tuple $\langle H, cs \rangle$, where H is a non-empty set of agents and cs a context set.

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- (6) Where p is a proposition and cs a context set,
 $cs \oplus p = cs \cap p$
- (7) **Assertion v. 1:** $c + \text{Assert}_a(\phi) = \langle H_c, cs_c \oplus \llbracket \phi \rrbracket \rangle$
Felicity condition in w : $\forall w' \in \text{Dox}_w(a) : w' \in \llbracket \phi \rrbracket$
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Felicity condition in w : $\forall w' \in \text{Dox}_w(a) : w' \in \llbracket \phi \rrbracket$
(‘a is committed to ϕ .’)
- (8) **Accommodating felicity inferences:** by default, if a move comes with a felicity condition f relative to w , as a precondition for interpreting that move in c , we take it that cs_c entails f .

Participation requirements

A general felicity condition:

- (9) A move α_a where a is some agent is felicitous in a context c only if $a \in H_c$.

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Some basic entrances and exits:

- (10) $c + \text{Enter}(a) = \langle H_c \cup \{a\}, cs_c \rangle$ (can be accommodated)

- (11) $c + \text{Exit}(a) = \langle H_c - \{a\}, cs_c \rangle$

Linguistic correlates? Leave this a question for now. (Cf. discussion in situated dialogue course.)

An example

The usual kind of thing: it's raining in w_1, w_2 and not in w_3, w_4 .

$\llbracket \text{it's raining} \rrbracket = \{w_1, w_2\}$

- (12) The scenario: a windowless room. A comes in from the outside.

$$c = \langle \{A, B\}, \{w_1, w_2, w_3, w_4\} \rangle$$

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$$c' = c + \text{Assert}_A(\ulcorner \text{it's raining} \urcorner)$$

$$c' = \langle H_c, c s_c \oplus \{w_1, w_2\} \rangle$$

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Responding to assertions

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Bledin & Rawlins 2016a,b)
B': What if Joanna is there?

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B': What if Joanna is there?
B': are you sure?
B': why not?

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- An assertion is a proposal to update the common ground with its content.
- In proposing, that assertion is put **on the 'table'**.

Tables

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- Farkas & Bruce (2010): intermediate stage between proposing an update to the common ground, and accepting it.
- An assertion is a proposal to update the common ground with its content.
- In proposing, that assertion is put **on the 'table'**.

Intuition: if an assertion is on the table, interlocutors are **coordinating on whether to incorporate it into the common ground.**

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(15) **Tabular assertion**

$c + \text{Assert}_a(\phi) = \langle H_c, \text{push}(A_c, \phi), cs_c \rangle$

Felicity condition in w : $\forall w' \in \text{Dox}_w(a) : w' \in \llbracket \phi \rrbracket$

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(16) **Acceptance**

$c + \text{Accept}_a = \langle H_c, \text{pop}(A_c), cs_c \oplus \llbracket \text{pop}(A) \rrbracket \rangle$

Felicity condition in w : $\forall w' \in \text{Dox}_w(a) : w' \in \llbracket \text{top}(A) \rrbracket$

(17) **Rejection**

$c + \text{Reject}_a = \langle H_c, \text{pop}(A_c), cs_c \rangle$

Felicity condition in w : $\forall w' \in \text{Dox}_w(a) : w' \notin \llbracket \text{top}(A) \rrbracket$

Example: acceptance of assertions

it's raining in w_1, w_2 and not in w_3, w_4 . $\llbracket \text{it's raining} \rrbracket = \{w_1, w_2\}$

- (18) The scenario: a windowless room. A comes in from the outside.

$$c_0 = \langle \{a, b\}, \langle \rangle, \{w_1, w_2, w_3, w_4\} \rangle$$

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For assertions, acceptance is the default! $c + \text{Assert}(\phi) + \text{Accept}$ amounts to assertion in v. 1.

Non-acceptance moves

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 - Resistance involves, at some level, a **strategy of inquiry** for deciding whether to accept an assertion.
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 - Resistance involves, at some level, a **strategy of inquiry** for deciding whether to accept an assertion.
 - Initial assertion remains on table while resistance move is dealt with.
- Assertion sequences (without acceptance) are mostly unconstrained so far. One more interesting case: sequences of contradictory assertions.

A Stalnakerian account for questioning

So far, we have only a single kind of inquiry: coordinating on a specific assertion.

- How can this be generalized?
- Starting point: modify the original Stalnakerian approach, and then return to a tabular approach.

Inquiry in a Stalnakerian context

We need a representation that can handle both **information** and **issues**.

- Information: what worlds are present at all.
- Issues: how do the worlds that are present relate to each other?

Groenendijk's 1999 idea: an equivalence relation on a subset of \mathcal{W} accomplishes this. (This leads to the notion of a **hybrid** in later work.)

A pre-formalizing example

Our usual four worlds. It's raining (only) in w_1, w_2 and snowing (only) in w_4 .

$$(19) \quad c + \text{「is it raining?」} = \left\{ \begin{array}{ll} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle, \\ & \langle w_3, w_3 \rangle, \quad \langle w_3, w_4 \rangle, \\ & \langle w_4, w_3 \rangle, \quad \langle w_4, w_4 \rangle \end{array} \right\}$$

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Intuition: cells correspond to ways the (informative) context set could evolve.

Another pre-formalizing example

$$(20) \quad c + \ulcorner \text{It's not snowing, but is it raining?} \urcorner =$$
$$\left\{ \begin{array}{ll} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle, \\ & \langle w_3, w_3 \rangle, \end{array} \right\}$$

Full update has eliminated one world altogether (w_4) and divided up w_1, w_2 from w_4 .

After Groenendijk (1999) (see Isaacs & Rawlins 2008, Ciardelli et al. 2013 etc. for descendants):

- (21) A **G-context set** is a set of pairs of worlds in some $cs \subseteq \mathcal{W}$ that is reflexive, symmetric, and transitive. (An equivalence relation.)
- (22) **Contexts v. 2**: A context is a tuple $\langle H, cs \rangle$, where H is a non-empty set of agents and cs a G-context set.

Inquiry in a Stalnakerian context: setup

Some convenience functions:

(23) Where Q is an equivalence relation:

- a. $\text{Dom}(Q) = \{w \mid \langle w, w \rangle \in Q\}$
- b. $\text{Alts}(Q) = \{p_{\langle st \rangle} \mid p \neq \emptyset \wedge \exists u_S : \forall v_S : \langle u, v \rangle \in Q \leftrightarrow p(v)\}$
- c. A proposition p **resolves** an equivalence relation Q iff $\exists p' \in \text{Alts}(Q) : p \subseteq p'$.¹

¹This is different than a Roberts-style complete answer.

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- A proposition p **resolves** an equivalence relation Q iff $\exists p' \in \text{Alts}(Q) : p \subseteq p'$.¹

Example: $\{w_1\}$ resolves $\left\{ \begin{array}{ll} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle, \\ & \langle w_3, w_3 \rangle \end{array} \right\}$

¹This is different than a Roberts-style complete answer.

Inquiry in a Stalnakerian context: moves

On to defining the moves. First, redefine \oplus , \otimes (here cf. Isaacs & Rawlins 2008).

(24) Where p is a proposition and c a context,
 $c \oplus p = c \cap \{\langle w, v \rangle \mid w, v \in p\}$

(25) Where p is a proposition and c a context,
 $c \otimes p = c \cap \{\langle w, v \rangle \mid w \in p \leftrightarrow v \in p\}$

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 $c \otimes p = c \cap \{\langle w, v \rangle \mid w \in p \leftrightarrow v \in p\}$

(26) **Assertion v. 2:** $c + \text{Assert}_a \phi = \langle H_c, cs_c \oplus \llbracket \phi \rrbracket \rangle$
Felicity conditions: the same (a is committed to ϕ)

(27) **Polar questions v. 1:** $c' = c + \text{PolarQ}_a \phi = \langle H_c, cs_c \otimes \llbracket \phi \rrbracket \rangle$
Felicity conditions in w : It is not the case that $\text{Dox}_a(w) \cap \text{Dom}(cs_{c'})$ resolves $cs_{c'}$.

A post-formalization example 1

Initial context

$$c_0 = \langle \{A, B\}, \left\{ \begin{array}{cccc} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, & \langle w_1, w_3 \rangle, & \langle w_1, w_4 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle, & \langle w_2, w_3 \rangle, & \langle w_2, w_4 \rangle, \\ \langle w_3, w_1 \rangle, & \langle w_3, w_2 \rangle, & \langle w_3, w_3 \rangle, & \langle w_3, w_4 \rangle, \\ \langle w_4, w_1 \rangle, & \langle w_4, w_2 \rangle, & \langle w_4, w_3 \rangle, & \langle w_4, w_4 \rangle \end{array} \right\} \rangle$$

Facts: it's raining (only) in w_1, w_2 and snowing (only) in w_4 .

A post-formalization example 2

Is it raining?

$$c_1 = \langle \{A, B\}, \left\{ \begin{array}{cccc} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, & \langle w_1, w_3 \rangle, & \langle w_1, w_4 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle, & \langle w_2, w_3 \rangle, & \langle w_2, w_4 \rangle, \\ \langle w_3, w_1 \rangle, & \langle w_3, w_2 \rangle, & \langle w_3, w_3 \rangle, & \langle w_3, w_4 \rangle, \\ \langle w_4, w_1 \rangle, & \langle w_4, w_2 \rangle, & \langle w_4, w_3 \rangle, & \langle w_4, w_4 \rangle \end{array} \right\} \otimes \{w_1, w_2\} \rangle$$

$$c_1 = c_0 + \ulcorner \text{is it raining?} \urcorner = \langle H_c, CS_c \otimes \llbracket \text{it is raining} \rrbracket \rangle$$

A post-formalization example 3

A: Is it raining?

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$$c_1 = c_0 + \ulcorner \text{is it raining?} \urcorner = \langle H_c, CS_c \otimes \llbracket \text{it is raining} \rrbracket \rangle$$

A post-formalization example 4

B: Yes, it's raining.

$$c_1 = \langle \{A, B\}, \left\{ \begin{array}{ll} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle \end{array} \right\} \rangle$$

$$c_2 = c_1 + \lceil \text{It's raining} \rceil = \langle H_{c_1}, CS_{c_1} \oplus \llbracket \text{it is raining} \rrbracket \rangle$$

A post-formalization example 4

B: Yes, it's raining.

$$c_1 = \langle \{A, B\}, \left\{ \begin{array}{ll} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle \end{array} \right\} \rangle$$

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- The context is now **uninquisitive**.

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B: Yes, it's raining.

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- The context is now **uninquisitive**.
- Relevance constraint after Roberts:

(28) A question-response α is **relevant** in a G-context c just in case there is some $p \in \text{Alts}(CS_c)$ such that $\llbracket \alpha \rrbracket$ decides p or $\llbracket \alpha \rrbracket$ decides $\neg p$.

Moving to non-polar questions

How to get from polar to constituent questions? (Here I diverge quite a bit from Groenendijk.)

- Intuition: can get the effect of a constituent question with a set of polar questions of this type.

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How to get from polar to constituent questions? (Here I diverge quite a bit from Groenendijk.)

- Intuition: can get the effect of a constituent question with a set of polar questions of this type.
- ‘What is the weather like?’ ~ ‘is it raining?’ + ‘is it sunny?’ + ‘is it snowing?’
- Suppose that a question denotation in general is a Hamblin alternative set (assume **mutual exclusivity** and **exhaustivity**).

Foreshadowing the details of polar questions

This generalizes the starting analysis of polar questions **as long as polar questions denote singleton sets.**

Restrict to alternative sets that partition some subset of \mathcal{W} (no overlap).

(29) Where Q is an alternative set and c a context,
 $c \oplus p = c \cap \{\langle w, v \rangle \mid \forall p \in Q : w \in p \leftrightarrow v \in p\}$

(30) **Questions v. 2.1**

$$c' = c + \text{Question}_a \phi = \langle H_c, \cap \{cs_c \otimes p \mid p \in \llbracket \phi \rrbracket\} \rangle$$

Felicity conditions in w : It is not the case that $\text{Dox}_a(w) \cap \text{Dom}(cs_{c'})$ resolves $cs_{c'}$.

Suppose it's raining in w_1, w_2 , sunny in w_3 and snowing in w_4 .

$\llbracket \text{What's the weather like?} \rrbracket = \{\{w_1, w_2\}, \{w_3\}, \{w_4\}\}$.

$\cap \{cs_c \circ p \mid p \in \llbracket \text{what's the weather like} \rrbracket\} =$

$$\left\{ \begin{array}{ll} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle, \\ & \langle w_3, w_3 \rangle, & \langle w_3, w_4 \rangle, \\ & \langle w_4, w_3 \rangle, & \langle w_4, w_4 \rangle \end{array} \right\}$$

$$\cap$$

$$\left\{ \begin{array}{lll} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, & \langle w_1, w_4 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle, & \langle w_2, w_4 \rangle, \\ & \langle w_3, w_3 \rangle, & \\ \langle w_4, w_1 \rangle, & \langle w_4, w_2 \rangle, & \langle w_4, w_4 \rangle \end{array} \right\}$$

$$\cap$$

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Questions and the table

Integrating tables, first pass

Can simply add an assertion stack to the G-context structure.
Is this enough?

How to incorporate tables into this picture?

- **assertions**: coordinating on evolution of the common ground.
 - Interaction with content: acceptance.
 - **Common ground management** (Repp 2013): rejection, postponement (others).

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 - Common ground management?

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- **assertions**: coordinating on evolution of the common ground.
 - Interaction with content: acceptance.
 - **Common ground management** (Repp 2013): rejection, postponement (others).
- **questions**: coordinating on goals of an inquiry.
 - Interaction with content: (partially) resolve.
 - Common ground management? reject question, start subinquiry, clarify, ...

(31) Contexts v. 3

A context is a tuple $\langle H, Q, A, cs \rangle$, where H is a non-empty set of agents, Q and A are stacks of sentences, and cs is a (regular) context set.

(32) Tabular assertion v. 2 (additional felicity conditions to be filled in)

$c + \text{Assert}_a(\phi) = \langle H_c, \text{push}(A_c, \phi), Q_c, cs_c \rangle$

Felicity condition in w : $\forall w' \in \text{Dox}_w(a) : w' \in \llbracket \phi \rrbracket$

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(33) Acceptance v. 2

$c + \text{Accept}_a = \langle H_c, \text{pop}(A_c), Q_c, cs_c \oplus \llbracket \text{pop}(A) \rrbracket \rangle$

Felicity condition in w : $\forall w' \in \text{Dox}_w(a) : w' \in \llbracket \text{pop}(A) \rrbracket$

Contexts with tables (2)

(34) Where p is a proposition, $inq(p) = \{\langle w, v \rangle \mid w, v \in p\}$

(35) **The i:** where c is a context,

$$QUD(c) = \begin{cases} \bigcap \{inq(cs_c) \cap p \mid p \in \llbracket \mathbf{top}(Q_c) \rrbracket\} & \text{if } |Q_c| \geq 1 \\ inq(cs_c) & \text{otherwise} \end{cases}$$

(36) **Dispelling a question:** where c is a context,

$$c + \text{Dispel} = \langle H_c, A_c, \text{pop}(Q_c), cs_c \rangle \text{ Felicitous only if } |Q_c| \geq 1$$

(37) **The full QUD in a context:** where c is a context,

$$FQUD(c) = \begin{cases} inq(cs_c) & \text{if } |Q_c| = 0 \\ QUD(c) \cap FQUD(c + \text{Dispel}) & \text{otherwise} \end{cases}$$

Contexts with tables (2)

(38) Questions with the table

$$c' = c + \text{Question}_a(\phi) = \langle H_c, \text{push}(Q_c, \phi), A_c, cs_c \rangle$$

Felicity conditions: appropriate in c at w only if

(i) If $|Q_c| \geq 1$ then $FQUD(c) \subseteq QUD(c')$, and

(ii) It is not the case that $\text{Dox}_a(w) \cap cs_{c'}$ resolves $QUD(c')$.

(39) Automatic dispelling

At any point c_n in a conversation, if $QUD(c_n) = \text{inq}(cs_{c_n})$,
adjust c_n to $c'_n = c_n + \text{Dispel}$.

Contexts with tables (3)

Relevance again:

- (40) A question-response α is **relevant** in a table context c just in case $\text{Alts}(QUD(c + \llbracket \alpha \rrbracket)) \subset \text{Alts}(QUD(c))^2$

²This is still different from Roberts-style relevance.

Polar questions again

Current analysis of the semantics of polar questions is a departure from Hamblin:

(41) $\llbracket \text{Is it raining?} \rrbracket = \{\lambda w_s. \text{it's raining in } w\}$

How to think about question-question sequences?

(42) What's the weather like? Is it raining?

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How to think about question-question sequences?

(42) What's the weather like? Is it raining?

In the G-context system, this would involve a redundant update. **But this seems felicitous!**

These are already licensed in the current system.

Polar questions again (2)

Licensing question-question sequences. Where c is the initial context:

$$QUD(c + \lceil \text{What's the weather like?} \rceil) = \left\{ \begin{array}{ll} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle, \\ & \langle w_3, w_3 \rangle, \\ & \langle w_4, w_4 \rangle \end{array} \right\}$$

is a subset of

$$QUD(c + \lceil \text{What's the weather like?} \rceil + \lceil \text{Is it raining?} \rceil) = \left\{ \begin{array}{ll} \langle w_1, w_1 \rangle, & \langle w_1, w_2 \rangle, \\ \langle w_2, w_1 \rangle, & \langle w_2, w_2 \rangle, \\ & \langle w_3, w_3 \rangle, & \langle w_3, w_4 \rangle, \\ & \langle w_4, w_3 \rangle, & \langle w_4, w_4 \rangle \end{array} \right\}$$

Polar questions again (3)

(43) Where should we go for lunch? Should we go to Mamoun's?

Biezma & Rawlins (2012): the function of a polar question relative to a bigger QUD is to characterize an alternative by 'name' – **identify constraint on the domain**.

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- The felicity condition acts as an **informative presupposition** (Prince 1978, Stalnaker 1973, 1974, a.o.)
- Biezma & Rawlins (2012) suggest that polar questions can **never** establish a big question. Stronger than the present constraint: could implement by adding a polar-specific presupposition (content alternative is part of the input QUD).

Alternative questions

Similar puzzle arises for alternative questions. On a naive implementation in a G-context system, they would involve redundant updates:

- (44) Where should we go for lunch? Should we go to Mamoun's or to Tacoria? (falling pitch)

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Biezma & Rawlins (2012) proposal – alternative questions list by 'name' all of the propositions in the current QUD. Implicate falling pitch in this (though this is controversial; see ?). Sketch:

- (45) Where α is a disjunction structure, $[[\alpha + \text{falling pitch}]]^c = [[\alpha]]^c$
Presupposes: $QUD(c) = QUD(c + [[\alpha]]^c)$

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Presupposes: $QUD(c) = QUD(c + [[\alpha]]^c)$

- This may force accommodation that eliminates alternatives that are in principle viable in c .

Summary

What have we accomplished?

- Core answers. (Fairly standard machinery in an update semantics context.)
- Basics of rejections / dismissals for assertions and questions.
- Room for resistance, strategies for acceptance – but not the full story.
- Question-question sequences and subquestions.

What's still missing?

- Weak answers (possibility claims, ignorance claims).
- Presupposition denials.
- A fuller story for resistance. (Probably not this class.)

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