

Geometric Langlands for ℓ -adic sheaves

Everything today is joint with Arinkin,
Gaitsgory, Kazhdan, Rozenblyum, and Varshavsky.

Background: $X/\mathbb{F}_q$ is a smooth, proj.
geom.-connected alg. curve. G split reductive,
 $\check{G}$ = dual group.

Langlands conj.

$$\left\{ \begin{array}{l} \text{Hecke eigenfunctions} \\ \text{on } G(\mathbb{A})/G(\mathcal{O}) \\ G(\mathbb{F}_q(X)) \end{array} \right\} \xleftrightarrow{\sim} \left\{ \begin{array}{l} \text{reps of} \\ \text{Weil}(X) \end{array} \right\}.$$

$\parallel$
 $Bun_G(\mathbb{F}_q)$

$\downarrow$
 $\check{G}(\mathbb{Q}_\ell)$

Bun_G = moduli stack of
 G -bundles on X .

Various geometric versions.

One, due to Beilinson-Drinfeld
(on this form, Arinkin-Gaitsgory):

Conj: X/k char $k=0$

$$\Rightarrow D\text{-mod}(Bun_G) \simeq \text{IndCal}_{N:1p}(LS_{\check{G}})$$

compatibly with Hecke functors.

Here: $LS_{\check{G}}$ is the ^(derived) moduli stack
of $\check{G}$ -bundles on X with ∇ .

C.F.:

Arithmetic Langlands vs Geom.
Langlands
SS SS

$$\{S' \longrightarrow S\} \simeq \mathbb{A}$$

$$L^2(S') \simeq L^2(\mathbb{A}).$$

Problem: char $k = 0 \stackrel{?}{\Rightarrow}$ no relevance to actual automorphic forms.

First issue: $LS_{\check{G}}$ is very specialized to dR setting.

Example: $\check{G} = G_m$:

$$0 \rightarrow \Gamma(X, \Omega^1) \rightarrow LS_{G_m} \rightarrow \text{Jac}(X) \times (BG_m)^{\times} \rightarrow 0$$

Main idea: introduce a moduli space $LS_{\tilde{G}}^{\text{restr}}$ in the ℓ -adic setting suitable for our purposes.

Notation: $k = \overline{k}$, X/k sm proj. connected

$\tilde{G}/\overline{\mathbb{Q}}_{\ell}$, G/k , $\ell \in k^{\times}$.

$\text{Lisse}(X) := (\text{DG})$ category of ^{ind-}lisse ℓ -adic sheaves on X .

Def: S affine test scheme, a

map $S \longrightarrow LS_{\tilde{G}}^{\text{restr}}$ is the

data of a symmetric monoidal

functor

$$\text{Rep } \check{Q} \longrightarrow \text{Lisse}(X) \otimes \mathcal{Q}\text{Coh}(S).$$

~~that~~ that is right t -exact.

Concretely: ~~$\mathbb{A}^1$~~ $S = \text{Spec } A$, $\text{RHS} = A$ -modules
in $\text{Lisse}(X)$.

Rem: If you replace $\text{Lisse}(X)$ w/ $\text{D-mod}(X)$,
get $LS_{\check{Q}}$ from before.

Examples: 1) $\check{Q} = \mathcal{Q}_m$:

$$LS_{\mathcal{Q}_m}^{\text{restr}} = \coprod_{\substack{\sigma \text{ rk } 1 \\ \text{lisse on } X}} \text{IB}\mathcal{Q}_m \times H'_{\text{ét}}(X, \bar{\mathbb{Q}}_l)_0 \times \Omega A'.$$

$\uparrow$
 derived
 line.

$$2) LS_{\mathcal{Q}_a}^{\text{restr}} = \text{IB}\mathcal{Q}_a \times H'_{\text{ét}}(X) \times \Omega A'.$$

$\nearrow$
 not formal.

3) $\check{G}$ gen'l : Then

$$LS_{\check{G}}^{\text{restr}} = \coprod_{(\check{M}, \sigma) \sim} LS_{\check{G}, \sigma}^{\text{restr}}$$

where $\check{M}$ is a Levi of $\check{G}$
and σ is a ~~semi-simple~~ irreducible
 $\check{M}$ -~~local~~ local system on X

ℓ $LS_{\check{G}, \sigma}^{\text{restr}}$ "looks like" the formal
completion of $LS_{\check{G}}$ along the
loc.-closed substack of local systems
w/ semi-simplification σ .



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$$\text{Conj: } \text{Shv}_{\text{Nilp}}(\text{Bun}_G) \simeq \text{IndCh}_{\text{Nilp}}(LS_{\check{G}}^{\text{vestr}})$$

compatibly w/ Hecke functors.

there: $\text{Shv}(\text{Bun}_G)$ is the category of (ind-constructible)

$\mathbb{Q}_\ell$ -sheaves on Bun_G , and Shv_{Nilp}

is the subcategory of sheaves w/

sing. support in $\text{Nilp} \subseteq T^* \text{Bun}_G$.



theory developed by Deligne, Beilinson, Saito.

in the ℓ -adic setting.

Idea to consider $\text{Shv}_{\text{Nilp}}(\text{Bun}_G)$ is
taken from the "Betti GLC" of
Ben-Zvi & Nadler.

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Thm: $\mathcal{Q}ch(LS_{\check{A}}^{\text{restr}}) \hookrightarrow \mathcal{Shv}_{\text{Nilp}}(\text{Bun}_q)$
 (compatibly w/ Hecke functors.)

Our proof is softer than Gaitsgory's
 analogous argument for D-modules.

$$1) \mathcal{C} = \mathcal{Shv}_{\text{Nilp}}(\text{Bun}_q).$$

Given $x \in X$ and $V \in \text{Rep } \check{A}$

$$\sim, H_{V,x}: \mathcal{Shv}_{\text{Nilp}}(\text{Bun}_q) \rightarrow \mathcal{Shv}_{\text{Nilp}}(\text{Bun}_q)$$

upgrade (due to Nadler-Yu.)

$$\text{Rep } \check{A} \otimes \mathcal{C} \longrightarrow \mathcal{C} \otimes \text{Lisse}(X).$$

variant: I finite set:

$$\text{Rep } \check{A}^I \otimes \mathcal{C} \longrightarrow \mathcal{C} \otimes \text{Lisse}(X^I).$$

2) Then: $\forall \mathcal{C}$. (mild tech condition)

this structure is $\Leftrightarrow \mathcal{Q}ch(LS_{\check{A}}^{\text{restr}}) \mathcal{R}\mathcal{C}$.

Reflects the Tannakian perspective on
 $LS_{\check{A}}^{\text{restr}}$.

Consequence: Self-duality of
 $\mathcal{Shv}_{N:lp}(Bun_G)$ ("regularized Verdier
duality for Bun_G ").

Take $K = \overline{\mathbb{F}}_q$, $X_0 / \overline{\mathbb{F}}_q$ descent of X/k .

Take $F: X \rightarrow X$ relative Frobenius.


$F: LS_{\tilde{a}}^{\text{restr}} \xrightarrow{\text{map}/\mathbb{Q}_\ell} LS_{\tilde{q}}^{\text{restr}}$ corresponding
 map ("pullback along $F: X \rightarrow X$ ").

Def. $LS_{\tilde{q}}^{\text{arithm}} := (LS_{\tilde{q}}^{\text{restr}})^{F=\text{id}}$

Think of $LS_{\tilde{q}}^{\text{arithm}}$ as a moduli space
 of Weil ~~rep's~~ group rep's.

Thm. $LS_{\tilde{q}}^{\text{arithm}}$ is a quasi-algebraic
 stack (not ind-, not infinite
 disjoint union).

Rem. Scholze and Zhu also introduced
 a version of this stack.

Example: $X_0 = \mathbb{P}^1 \Rightarrow LS_{\check{a}}^{a\text{-thm}} \cong \check{a}/\check{a}$.

Recall Hochschild construction: A alg

$M \in A\text{-mod}$ dualizable $A\text{-mod}$

$\sim [M] \in HH_*(A)$

Categorically: $\text{tr}(\text{id}_M \circ M) \in \text{tr}(\text{id}_{\text{Bun}_A})$.

In our setting:

$$\begin{array}{ccc}
 F^* & & F^* \\
 \downarrow & & \downarrow \\
 \text{QGr}(LS_{\check{a}}^{\text{restr}}) & \simeq & \text{Shv}_{\text{Nilp}}(\text{Bun}_A) \\
 \underbrace{\hspace{2cm}} & & \underbrace{\hspace{2cm}} \\
 A & & M
 \end{array}$$

$$\begin{aligned} \text{tr}(F \cap \mathcal{OCh}(LS_{\check{A}}^{\text{nestr}})) &\ni \text{tr}(F \cap \text{Shv}_{\text{Nilp}}(\text{Bun}_q)) \\ \parallel &\quad \parallel \\ \mathcal{OCh}(LS_{\check{A}}^{\text{arithm}}) &\ni \text{Drinf}. \end{aligned}$$

Thm-in-progress: $\Gamma(LS_{\check{A}}^{\text{arithm}}, \text{Drinf}) = \text{tr}(F \cap \text{Shv}_{\text{Nilp}}(\text{Bun}_q))$

$$\simeq \text{Fun}_c(\text{Bun}_q(\mathbb{F}_q)).$$

Rem: Uses Cong Xue's (following V. Lafforgue) work substantially.

Slogan:

automorphic forms localize onto

$LS_{\check{A}}^{\text{arithm}}$ (Drinfeld's picture).

Rem: Restricted GLC $\Rightarrow$

$$\text{Drinf} = \omega_{LS_{\check{A}}^{\text{arithm}}}.$$

Baby

Example: $X_0 = \mathbb{P}^1$. $LS_{\check{q}}^{\text{arithm}} = \bar{a}/\check{a}$

$$\Gamma(LS_{\check{q}}^{\text{arithm}}, \mathcal{O}) = \bigotimes \bar{\mathbb{Q}}_l[\check{\lambda}]^w.$$

vs. $\text{Bun}_q(\mathbb{F}_q) = \check{\lambda}^+$ ignoring stackyness

$$\bigotimes \text{Func}(\check{\lambda}^+) = \bigotimes \bar{\mathbb{Q}}_l[\check{\lambda}^+] = \bar{\mathbb{Q}}_l[\check{\lambda}]^w.$$